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A Category \mathcal{Q} consists of a collection of objects, denoted by $Ob(\mathcal{Q})$ and for each pair of objects $X, Y \in Ob(\mathcal{Q})$

a set $hom_{\mathcal{Q}}(X, Y) \neq \emptyset$ for each

object $X \in Ob(\mathcal{Q}) \ni$ a distinguished

morphism $1_X \in hom_{\mathcal{Q}}(X, X)$, called the identity morphism and for objects

$X, Y, Z \in Ob(\mathcal{Q})$, \exists a composition law

$hom_{\mathcal{Q}}(Y, Z) \times hom_{\mathcal{Q}}(X, Y) \longrightarrow hom_{\mathcal{Q}}(X, Z)$ which satisfies

The following axioms:

i) (identity) For 1_X , and 1_Y and a

morphism $f \in \text{hom}_\mathcal{C}(X, Y)$

$$1_Y \circ f = f = f \circ 1_X$$

(here ' \circ ' is
the composition
law).

We

denote a morphism

$$f \in \text{hom}_\mathcal{C}(X, Y)$$

by

$$f: X \rightarrow Y$$

\mathcal{C} Top

A subcategory \mathcal{C}' of a category \mathcal{C}

is a category \mathcal{C}' s.t

$$\text{Ob}(\mathcal{C}') \subset \text{Ob}(\mathcal{C})$$

and $\text{hom}_{\mathcal{C}'}(x', y') \subset \text{hom}_{\mathcal{C}}(x', y')$

$$x', y' \in \text{Ob}(\mathcal{C}') \subset \text{Ob}(\mathcal{C})$$

Example (i) Let $\mathcal{C} = \text{Grp}$ — the category of groups
 $\mathcal{S} = \text{Set}$ — the category of sets.

$$\text{Grp} \subset \mathcal{S}$$

Let $\mathcal{C} = \text{Grp}$ — category of groups

and $\mathcal{C}' = \text{Ab}$ — category of abelian groups

Let G, H be abelian groups $\text{hom}_{\text{Ab}}(G, H) = \text{hom}_{\text{Grp}}(G, H)$

A subcategory \mathcal{Q}' of \mathcal{Q} is called a full-subcategory if

for any objects $X, Y \in \text{Ob}(\mathcal{Q})$

$$\text{hom}_{\mathcal{Q}'}(X, Y) = \text{hom}_{\mathcal{Q}}(X, Y).$$

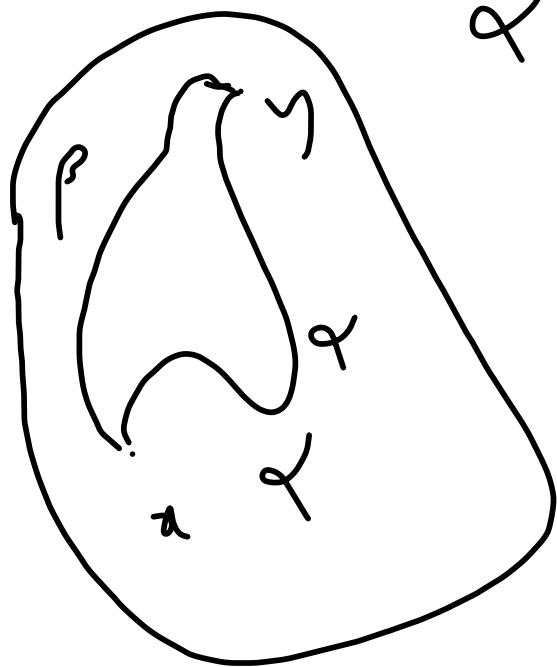
A category \mathcal{Q} is called small if $\text{Ob}(\mathcal{Q})$ is a set.

Example

Let X be a topological space.

Let $x, y \in X$. A path in X from x to y is a continuous function

$$\alpha: [0, 1] \rightarrow X \quad \text{s.t.} \quad \alpha(0) = x \\ \alpha(1) = y$$



Let $\beta: [0, 1] \rightarrow X$ be another path from x to y .

We say that " α is homotopic to β relative endpoints"

written as $\alpha \simeq \beta$ rel $(0,1)$

if \exists a continuous function

$$F: \underline{I} \times \underline{I} \longrightarrow X, \quad \underline{I} = [0,1]$$

Other notation:

For each $t \in [0,1]$

let $\alpha_t: [0,1] \rightarrow X$

be defined

by

$$\alpha_t(s) = F(s,t)$$

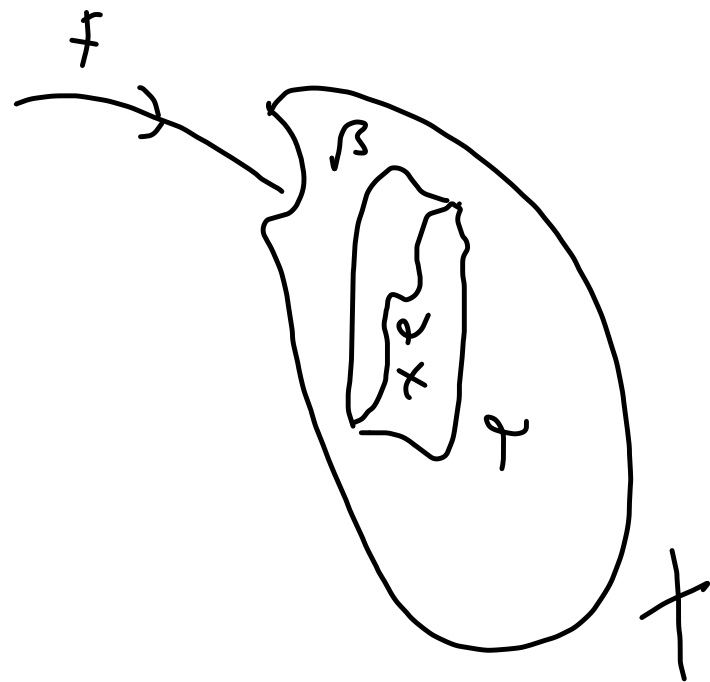
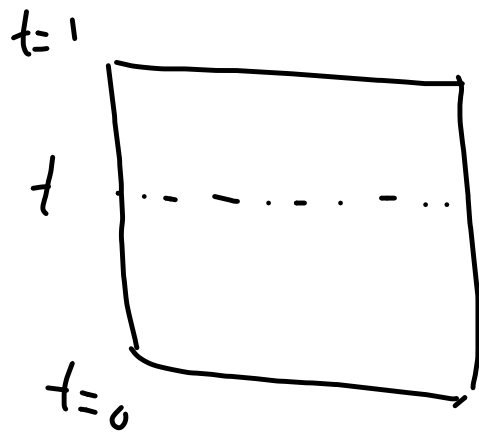
$$\alpha_t \quad \wedge \quad F(s,0) = \alpha(s)$$

$$F(s,1) = \beta(s)$$

$$F(0,t) = x$$

$$F(1,t) = y$$

$$\forall t \in [0,1]$$



Let $P(X; a, y)$ be the set of all
paths ^{in X} from x to y

Consider $P(X; a, y)$ as a topological space

Then Continuity (joint) of $F: \underline{I} \times \underline{I} \rightarrow X$
will ensure that the assignment

$$\begin{array}{ccc} \underline{I} & \longrightarrow & P(X; a, y) \\ t & \longmapsto & \gamma_t \end{array} \text{ is continuous.}$$

First note that " \sim rel $(0,1)$ "

is an equivalence relation on $P(X; z, y)$.

$[q]$ — equivalence class of q

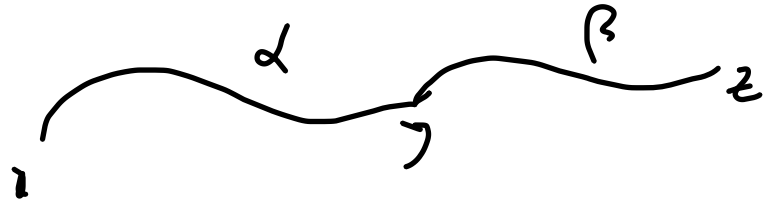
Given a space X ,
Define a category ΠX whose objects

are points of X

Given two points

$$\text{Hom}_{\Pi X}(x, y) = \left\{ [q] \mid \begin{array}{l} x, y \in X \quad (x, y \in \text{Ob}(\Pi X)) \\ \text{homotopy classes of paths} \\ \text{from } x \text{ to } y \end{array} \right\}$$

If $a, y, z \in X$



and $\alpha: a \rightarrow y$, $\beta: y \rightarrow z$ are paths

Then $\alpha * \beta: a \rightarrow z$

$$\alpha * \beta (s) = \begin{cases} \alpha(2s) & 0 \leq s \leq \frac{1}{2} \\ \beta(2s-1) & \frac{1}{2} \leq s \leq 1 \end{cases}$$

$$[\alpha] = \left\{ \alpha': [0,1] \rightarrow X \mid \begin{array}{l} \alpha'(0) = a \\ \alpha'(1) = y \end{array} \text{ and } \alpha' \simeq \gamma \text{ rel } (0,1) \right\}$$
$$[\alpha] \cdot [\beta] = [\alpha * \beta]$$

$$\frac{1}{2} \leq s \leq 1$$

Check that this is well-defined.

Funda
mental
Groupoid.

Small
Category.

$\alpha: x \rightarrow y$
 $[\alpha] := [\alpha^{-1}]$
 $\alpha^{-1}: y \rightarrow x$

$\alpha^{-1}(t) = \alpha(1-t)$

To check this is well defined
verify:

if $\alpha \simeq \alpha'$ rel(0,1), $\alpha, \alpha': x \rightarrow y$

and $\beta \simeq \beta'$ rel(0,1), $\beta, \beta': y \rightarrow z$

then $\alpha * \beta \simeq \alpha' * \beta'$ rel(0,1).

$= [c_x], c_x: I = [0,1] \rightarrow X$

$c_x(s) = \alpha \forall s$

Take $x=y$

$$P(X; x, x)$$

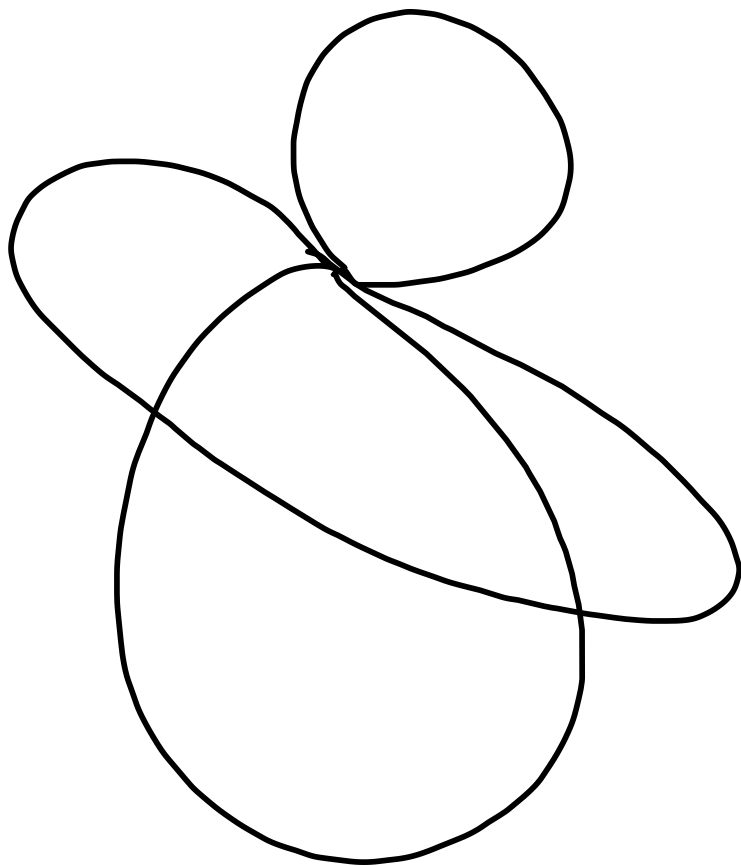
$$[X; x, x]$$

$$= P(x; x, x)$$

$$= \hat{\pi}_1(X, x) \cong \pi_1(X, x)$$

called the fundamental group of X .

A path from x to x
is a loop.



Functor: Let \mathcal{C} and \mathcal{D} be
two categories.

A Covariant functor $F: \mathcal{C} \rightarrow \mathcal{D}$
is an assignment which assigns
to each object $X \in \text{Ob}(\mathcal{C})$, an object
 $F(X) \in \text{Ob}(\mathcal{D})$ and to each
morphism $f: X \rightarrow Y$ in $\text{Hom}_{\mathcal{C}}(X, Y)$

a morphism

$$F(f): F(X) \rightarrow F(Y)$$

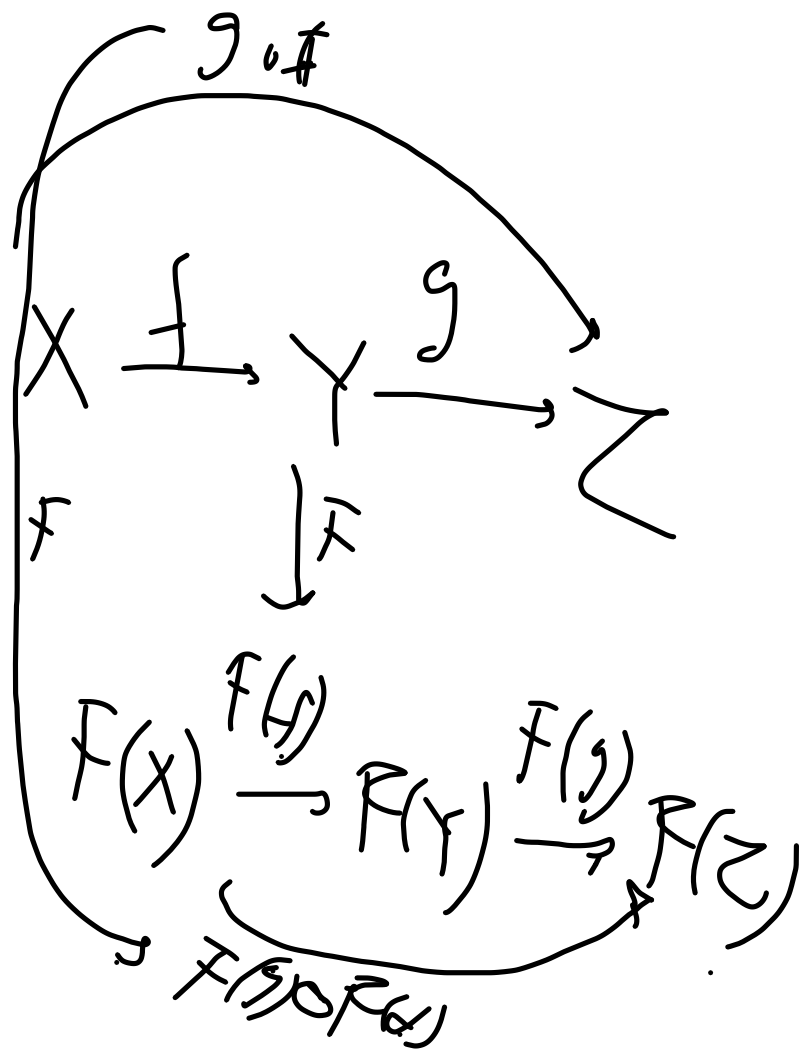
Similarly Contravariant
functor
is defined.

and the following rule must be

Satisfied.

$$(i) F(1_X) = 1_{F(X)}$$

$$(ii) F(g \circ f) = F(g) \circ F(f)$$



Let $F, G : \mathcal{C} \rightarrow \mathcal{D}$ be Covariant functors

A natural transformation

$$\eta : F \rightarrow G$$

is an assignment

$$X \in \text{Ob}(\mathcal{C}) \longrightarrow$$

$$f \downarrow \in \text{Hom}_{\mathcal{C}}(X, Y) \longrightarrow$$

$$Y \in \text{Ob}(\mathcal{C})$$

If $\eta(X)$ is an isomorphism $\forall X$
 then η is called an equivalence.
 Commutative.

$$\eta(X) : F(X) \rightarrow G(X)$$

$$\begin{array}{ccc} \downarrow F(f) & \circlearrowright & \downarrow G(f) \\ (Y) \in \text{Ob}(\mathcal{C}) & \xrightarrow{\eta(Y)} & (Y) \in \text{Ob}(\mathcal{C}) \end{array}$$