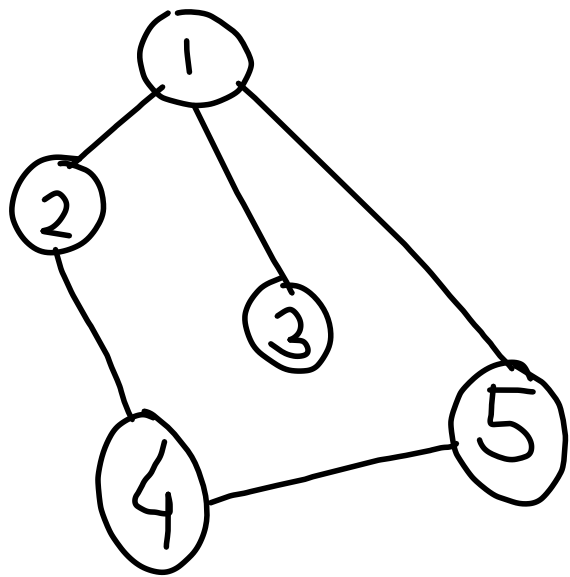


# Graph Algorithms

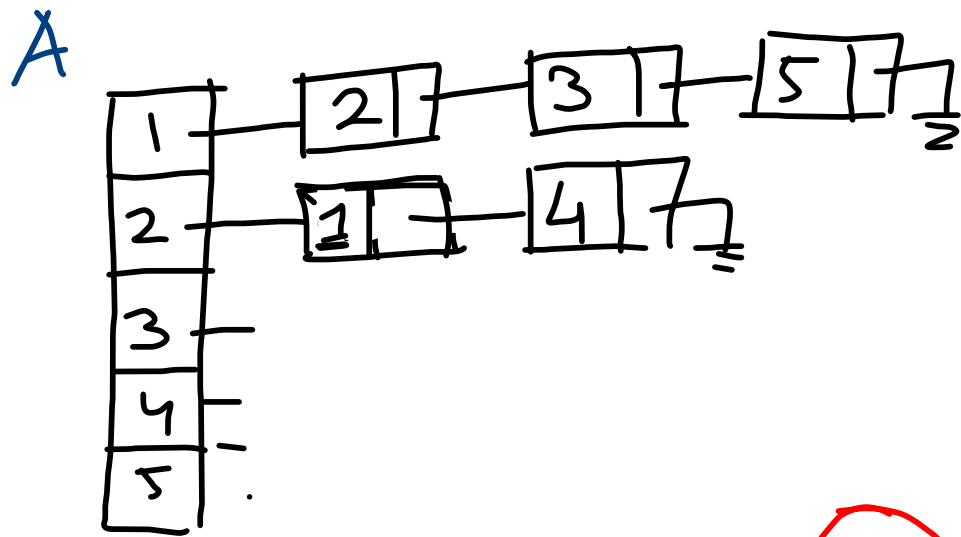
## Graph Representation:



$$G = \{ (1,2), (1,3), (1,5), (2,4), (4,5) \}$$

$$\text{Space} \rightarrow O(|E|)$$

Q Is there an edge between 3 and 5?  
 $O(|E|)$



$$\text{Space} = O(n + |E|)$$

Adjacency List

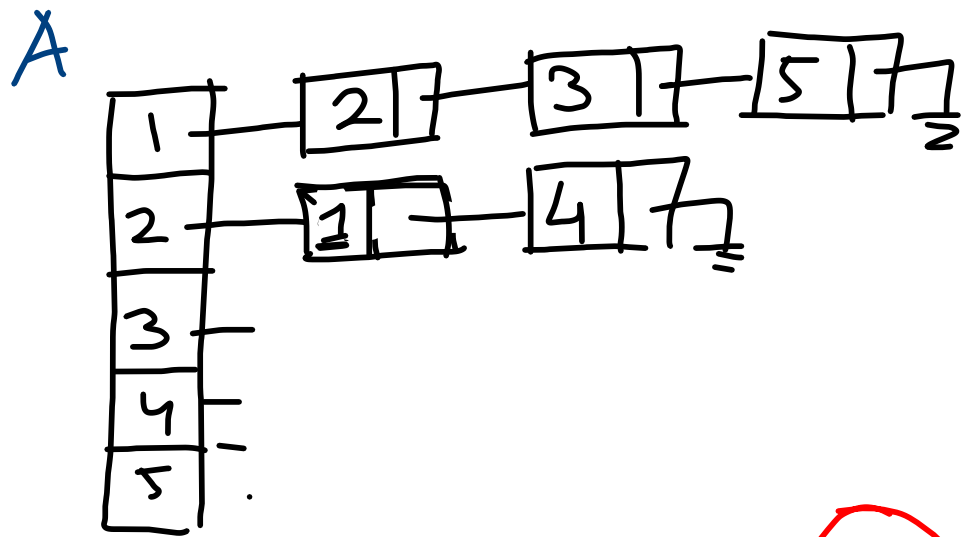
Q

$$O(\Delta)$$

↳ max degree

$$\approx O(n)$$

$$\underline{|V| = n}$$



$$\text{Space} = O(n + |E|)$$

Adjacency List

Q

$$O(\Delta)$$

↳ max degree

$$\approx O(n)$$

$$\underline{|V| = n}$$

# Adjacency Matrix

Space  $\rightarrow O(n^2)$

Q  $\rightarrow O(1)$

A

	1	2	3	4	5
1	0	1	1	0	1
2	-			-	
3					
4					
5		-			

$A[2,3] = ?$

1  $\rightarrow$  edge

0  $\rightarrow$  no edge

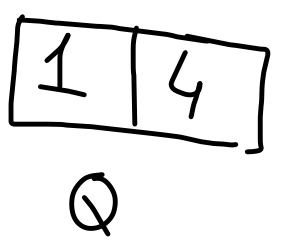
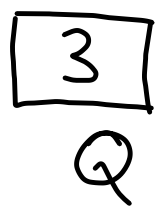
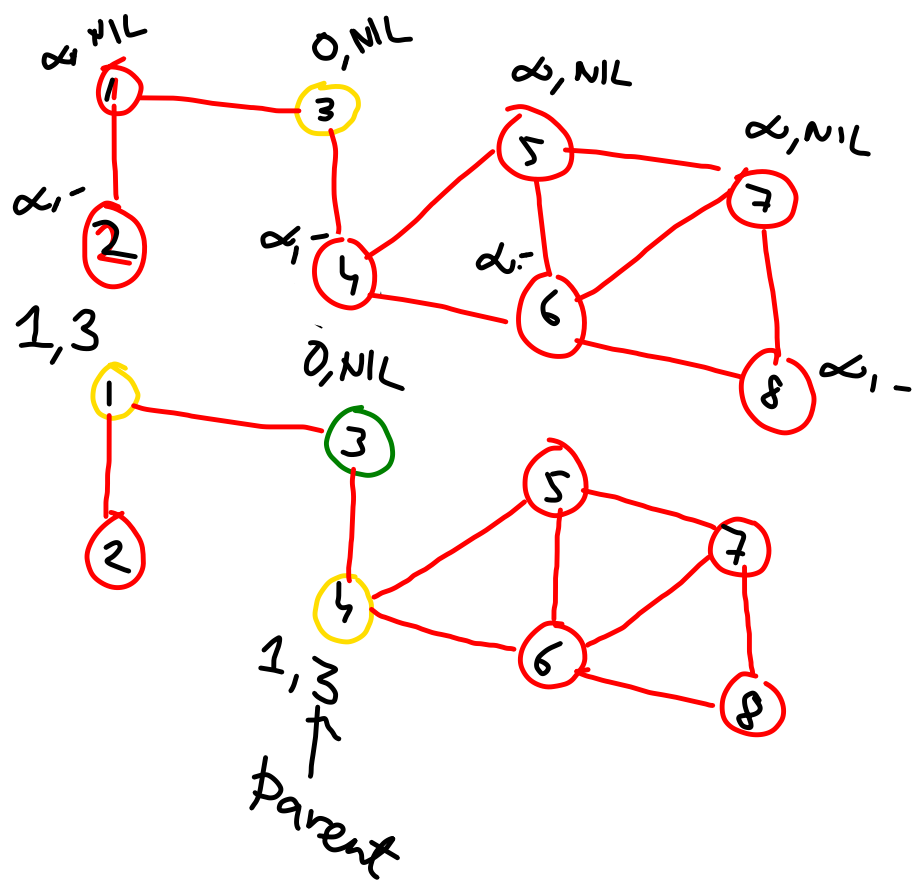
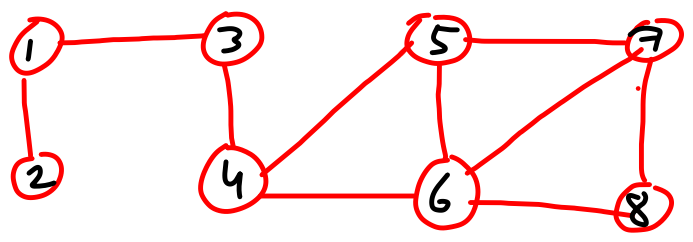
- Sparse  $\rightarrow$  Adjacency List
- Dense  $\rightarrow$  Adjacency Matrix

# Graph Search Algorithms

## Applications :

- Web crawling
- GPS navigation
- FB Friend finder
- N/W Broadcast
- Garbage Collection

# BFS (Breadth First Search)



## Search

- ↳ mindistance from  $s$  to  $t$
- ↳ path from  $s$  to  $t$  (mindist)
- ↳ Cycle detection

# BFS(G, s)

$\forall u \in V - \{s\}$

$u.color = RED$

$u.dist = \infty$

$u.parent = NIL$

$s.color = YELLOW$

$s.dist = 0$

$s.parent = NIL$

$Q = \emptyset$

Enqueue(Q, s)

Init

(H) (M+|E)

while ( $Q \neq \emptyset$ )

$u = DEQUEUE(Q)$

$\forall v \in Adj[u]$

if ( $v.color == RED$ )  
{

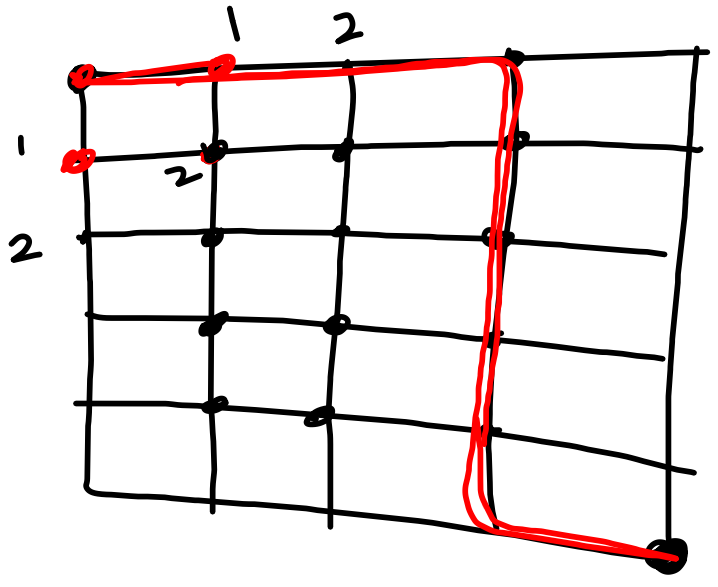
$v.color = yellow$

$v.dist = u.dist + 1$

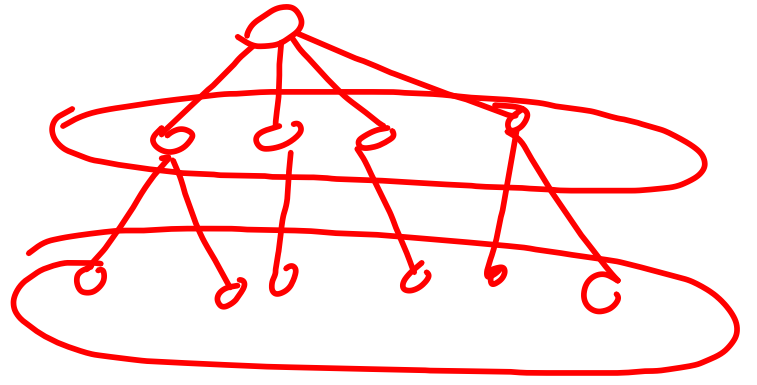
$v.parent = u$

Enqueue(Q, v)

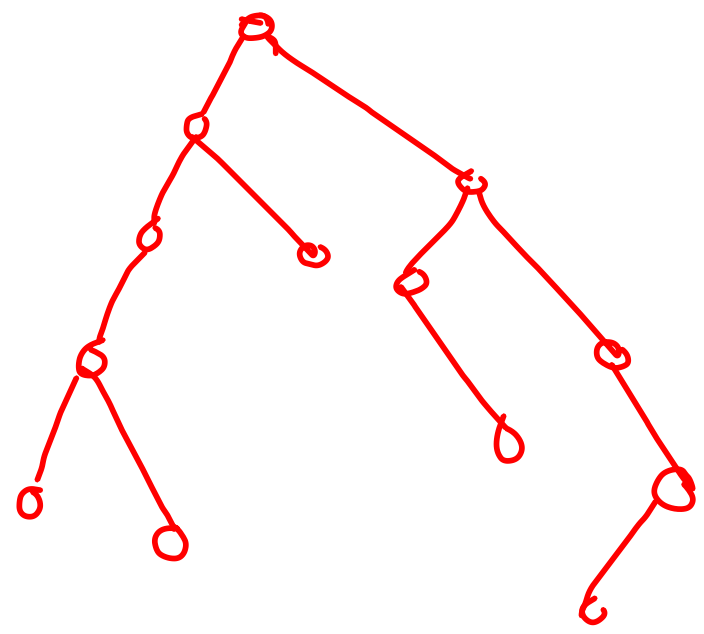
}  
 $u.color = GREEN$



BFS



⋮



DFS



# DFS (G)

$\forall u \in V$

$u.\text{color} = \text{RED}$

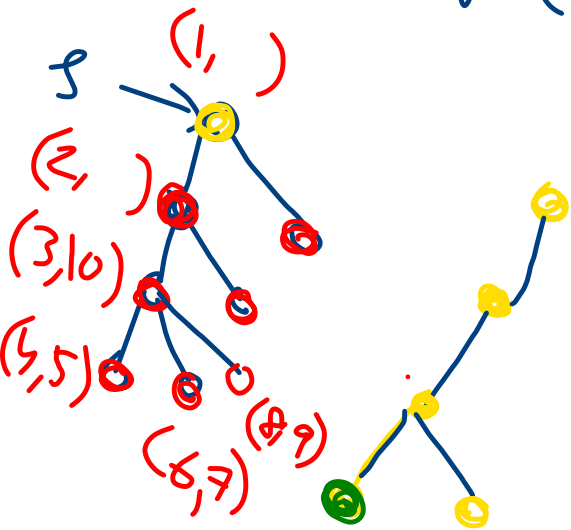
$u.\text{parent} = \text{NIL}$

$\text{time} = 0$

$\forall u \in V$

if ( $u.\text{color} == \text{RED}$ )

DFS\_VISIT(G, u)



# DFS\_VISIT(G, u)

$\text{time} = \text{time} + 1$

$u.\text{start} = \text{time}$

$u.\text{color} = \text{Yellow}$

$\forall v \in \text{Adj}[u]$

if ( $v.\text{color} == \text{RED}$ )

$v.\text{parent} = u$

DFS\_VISIT(G, v)

$u.\text{color} = \text{GREEN}$

$\text{time} = \text{time} + 1$

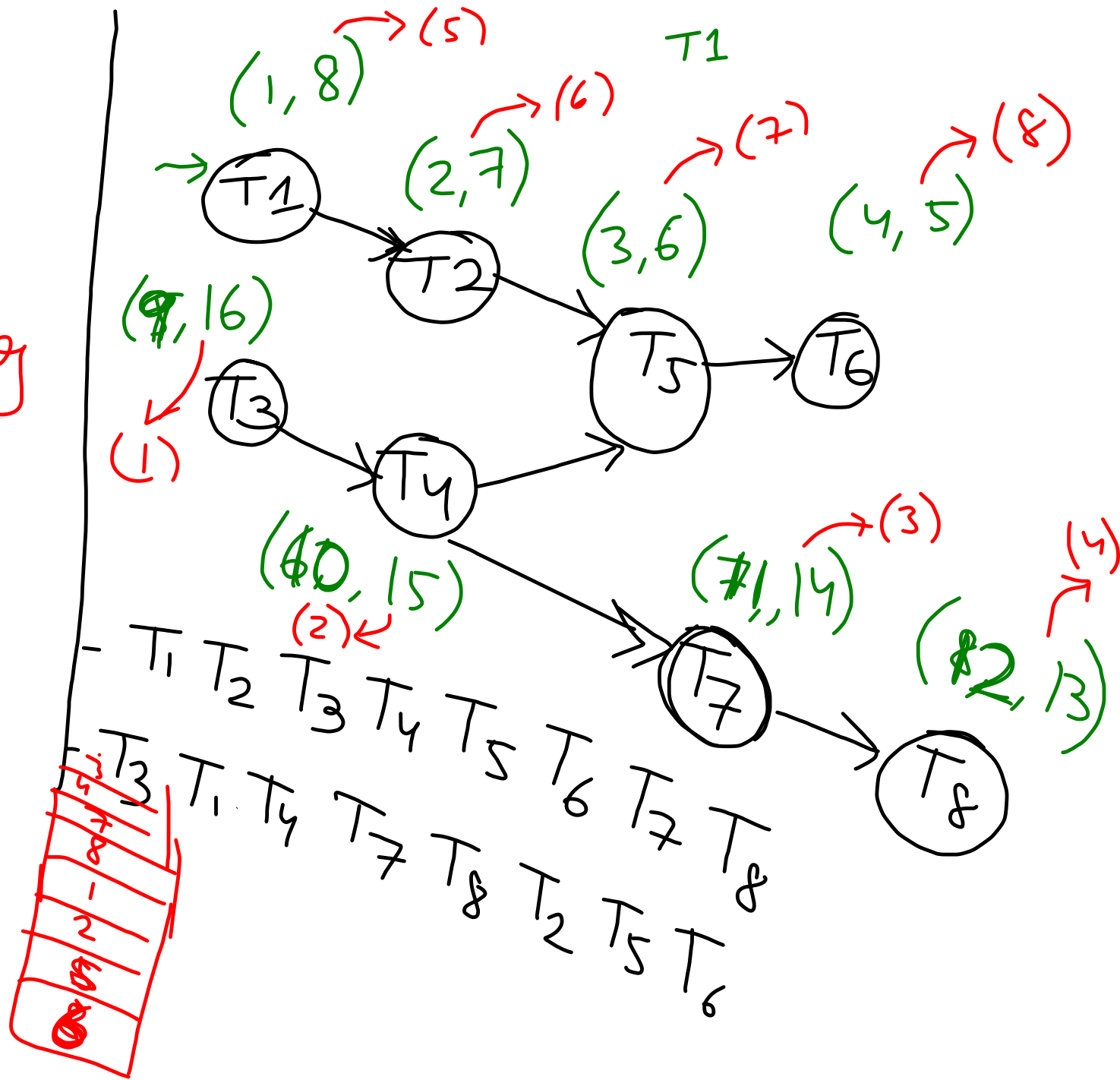
$u.\text{finish} = \text{time}$

# Topological Sort

Linear Ordering  
(One such ordering)

DAG

↓  
Direct Acyclic Graph



# Topological Sort

Linear Ordering  
(One such ordering)

DAG

↓  
Direct Acyclic Graph

