

# Assignment II: Symmetric Key Encryption

Deadline: **24/09/2022**

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- Let  $F$  be a secure PRF defined over  $(\{0, 1\}^n, \{0, 1\}^n, \{0, 1\}^n)$ .
  - Prove that  $G_k(x) := F_k(x) \oplus x$  is a secure PRF.
  - Prove that  $G_{(k,k')}(x) := F_k(x) \oplus F_{k'}(x)$  is a secure PRF.
  - Show that  $G_k(x) := F_k(x) \| F_k(F_k(x))$  is insecure.
- Prove that if  $F$  is a pseudorandom function defined over  $(\{0, 1\}^n, \{0, 1\}^n, \{0, 1\}^n)$ , then

$$G(s) := F_s(1) \| F_s(2) \| \cdots \| F_s(\ell)$$

is a pseudorandom generator with expansion factor  $n\ell$ .

- Let  $F$  be a pseudorandom function from  $n$ -bits to  $n$ -bits and  $G$  be a pseudorandom generator with expansion factor  $n + 1$ . For each of the following encryption schemes, state whether the scheme (i) has indistinguishable encryptions in the presence of an eavesdropper and whether (ii) it is IND-CPA secure.
  - $Enc_k(m) := \langle r, G(r) \oplus m \rangle$ , where  $m \in \{0, 1\}^{n+1}$ ,  $r \leftarrow_{\$} \{0, 1\}^n$ .
  - $Enc_k(m) := m \oplus F_k(0^n)$ .
  - $Enc_k(m) := \langle r, m_1 \oplus F_k(r), m_2 \oplus F_k(r + 1) \rangle$ , where  $r \leftarrow_{\$} \{0, 1\}^n$ , where  $m = m_1 \| m_2$ ,  $|m_1| = |m_2| = n$ .
- Let  $F$  be a pseudorandom permutation. Consider the mode of operation in which a uniform value  $r \leftarrow_{\$} \{0, 1\}^n$  is chosen, and the  $i$ -th ciphertext block  $c_i$  is computed as

$$c_i := F_k(r \oplus i \oplus m_i).$$

Show that this scheme does not have indistinguishable encryptions in the presence of an eavesdropper.

- Consider a ciphertext corresponding to a message encrypted using CBC mode. Suppose a single bit error has occurred in one of the ciphertext block. Which plaintext blocks will it affect during the decryption? What is the effect of such single-bit error in the ciphertext in case of Counter (CTR) mode?