


Grammars to PDA

$$G = (V, T, P, S)$$

PDA: $P' = (\{q\}, T, V \cup T, \underline{\delta}, q, S)$.

Empty
Stack



① \forall variable A ,

$$\delta(q, \epsilon, A) = \left\{ (q, \beta) \mid A \rightarrow \beta \text{ is a production in } P \right\}$$

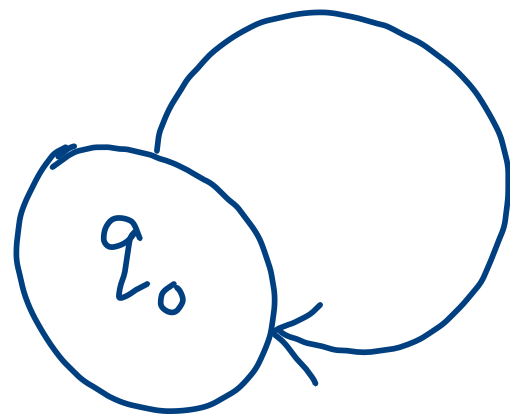
② \forall terminal a ,

$$\delta(q, a, a) = \{ (q, \epsilon) \}$$

Set of Palindrome's over $\{0,1\}$

$S \rightarrow 0S0 \mid 1S1 \mid 0 \mid 1 \mid \epsilon$

PDA



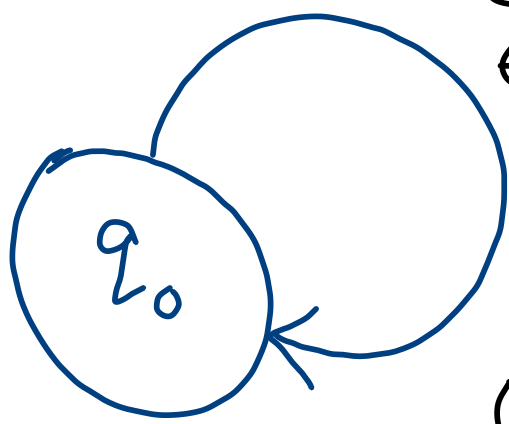
Set of Palindrome's over $\{0,1\}$

$S \rightarrow 0S0 \mid 1S1 \mid 0 \mid 1 \mid \epsilon$

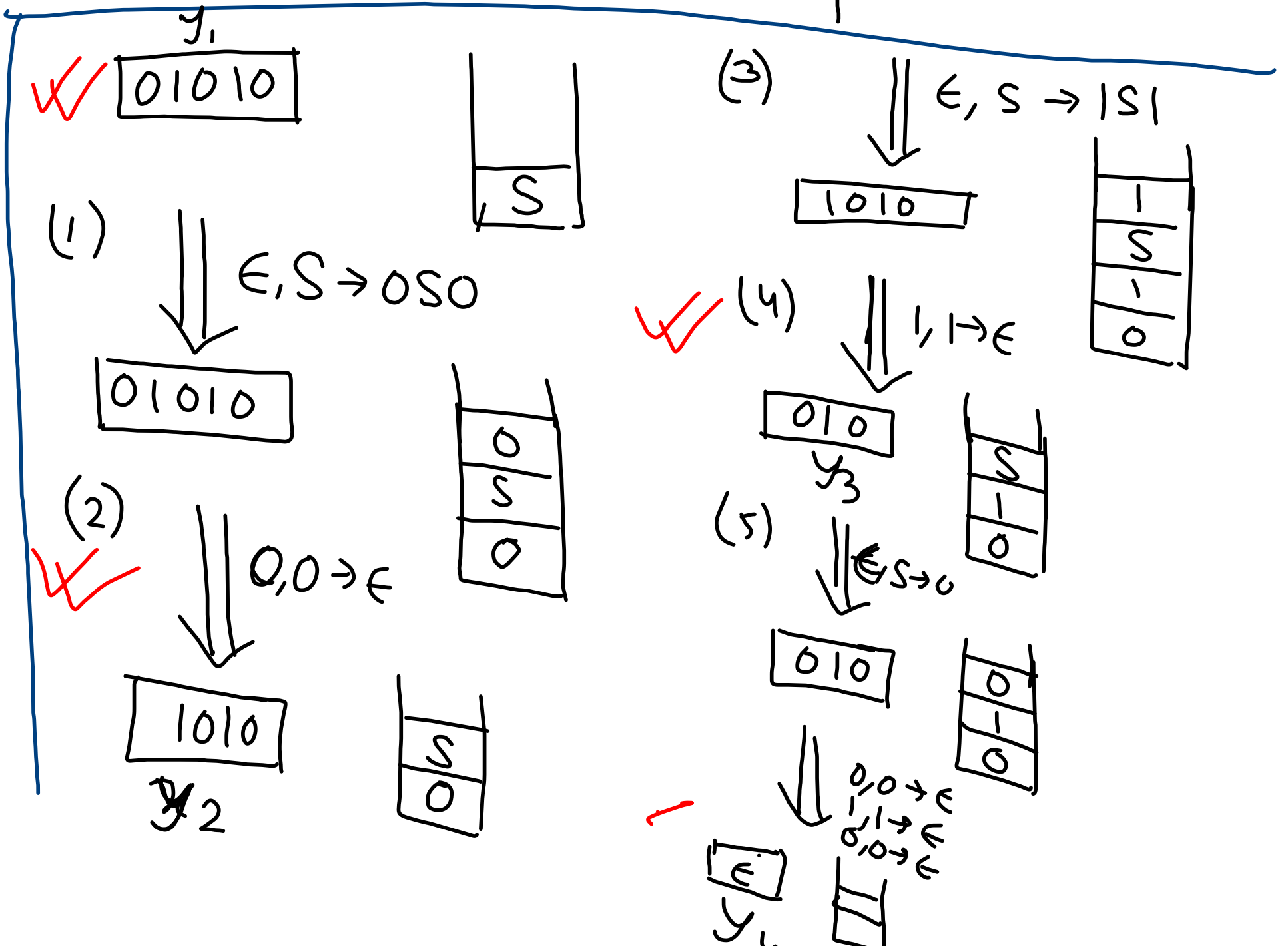
01010 (w)
 $S \rightarrow 0\underline{S}0$
 $\rightarrow 01S10$
 $\rightarrow 01010$

left most derivation

PDA



- $\epsilon, S \rightarrow 0S0$
- $\epsilon, S \rightarrow 1S1$
- $\epsilon, S \rightarrow 0$
- $\epsilon, S \rightarrow 1$
- $\epsilon, S \rightarrow \epsilon$
- $0, 0 \rightarrow \epsilon$
- $1, 1 \rightarrow \epsilon$
- $\epsilon, \epsilon \rightarrow \epsilon$



$$\underline{\omega \in L(G) \Rightarrow \omega \in N(P')}$$

Sentential Form

$$(q_0, \omega, S) \xrightarrow{*} (q_0, \epsilon, \epsilon)$$

⇓

{ leftmost derivation

⇓

$$\gamma_i = \underline{\alpha_i} \alpha_i$$



$\gamma_2 =$	$\frac{O}{\alpha_2}$	$\frac{SO}{\alpha_2}$
$\gamma_3 =$	$\frac{O1}{\alpha_3}$	$\frac{SIO}{\alpha_3}$
$\gamma_4 =$	$\frac{OI0IO}{\alpha_4}$	$\frac{\epsilon}{\alpha_4}$

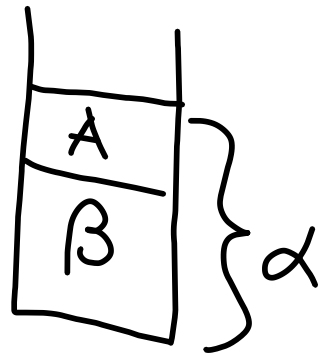
$$\leftarrow \alpha_i \mid \gamma_i = \omega$$

$$\underbrace{(q_0, w, S)}_{\gamma_1} \rightarrow \underbrace{(q_0, y_1, \alpha_1)}_{\gamma_2} \rightarrow \dots \rightarrow (q_0, \epsilon, \epsilon)$$

$$S \xRightarrow{*} \gamma_1 \xRightarrow{*} \gamma_2 \dots \xRightarrow{*} \gamma_m = w \leftarrow \text{CFG}$$

$$(q_0, y_i, \alpha_i) \rightarrow (q_0, \underline{y_{i+1}}, \alpha_{i+1}) \quad (\text{Induction})$$

y_i

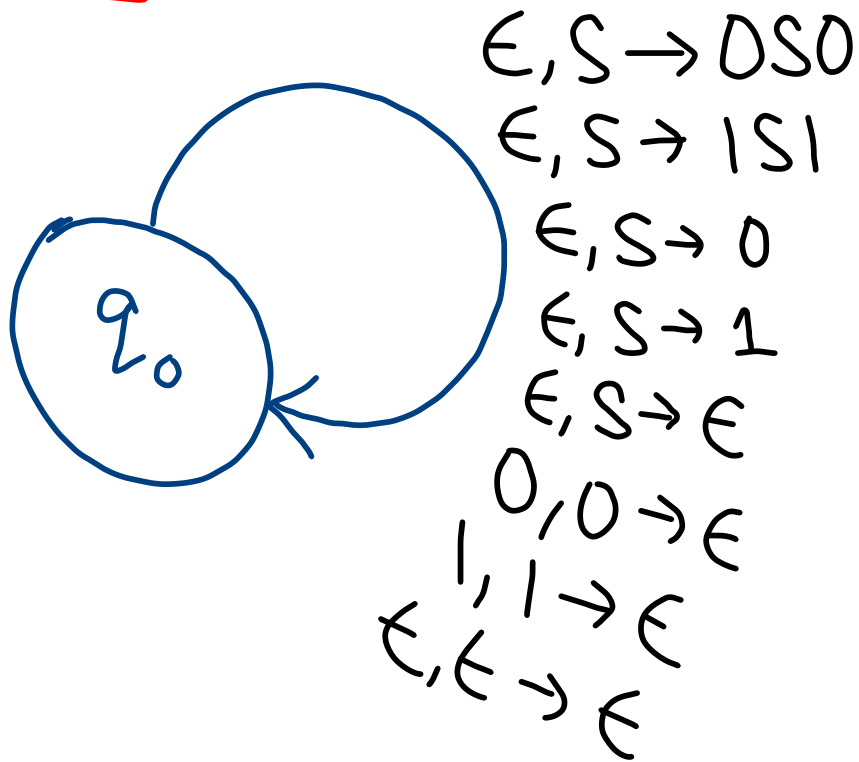


Set of Palindrome's over $\{0,1\}$

$$S \rightarrow OSO \mid 1S1 \mid 0 \mid 1 \mid \epsilon$$

01010
 $S \rightarrow 0\underline{S}0$
 $\rightarrow 01S10$
 $\rightarrow 01010$

PDA



$$E \rightarrow E * E \mid E + E \mid a \mid b$$

$$\underline{a * a + a}$$

$$E \rightarrow E * E$$

$$\rightarrow a * E + E$$

$$\rightarrow a * a + a$$

(left most derivation)

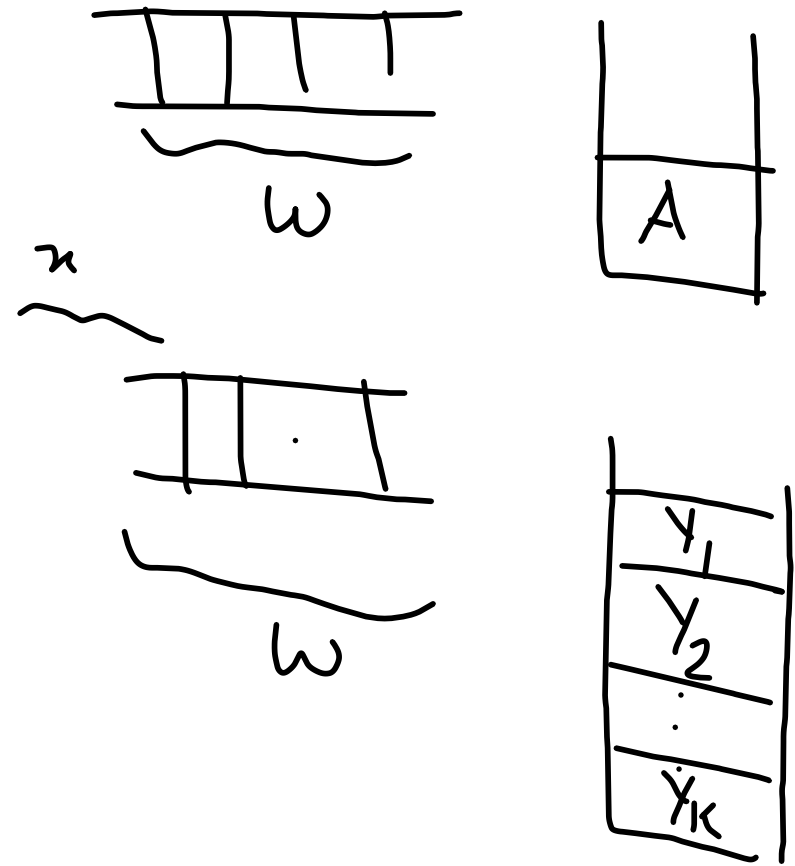
$$E \rightarrow E + E$$

$$\rightarrow E * E + a$$

$$\rightarrow a * a + a$$

$$\underline{w \in \mathcal{N}(P') \implies w \in \mathcal{L}(G)}$$

$$A \rightarrow \underbrace{y_1 y_2 \dots y_k}$$



Induction on k

$$y_i \stackrel{*}{\implies} x_i$$

$$A \stackrel{*}{\implies} y_1 y_2 \dots y_k$$

$$\Downarrow$$

$$x_1 x_2 \dots x_k = \underline{\underline{w}}$$

Closure Properties

① Union: $(L_1 \cup L_2)$

$L_1, L_2 \rightarrow \text{CFL}$

$L_1 \cup L_2 \rightarrow \text{CFL?}$

$$G_1 = (V_1, \Sigma, P_1, S_1)$$

$$G_2 = (V_2, \Sigma, P_2, S_2)$$

$$G = (V, \Sigma, P, S)$$

$$P = P_1 \cup P_2 \cup \{S \rightarrow S_1 \mid S_2\}$$

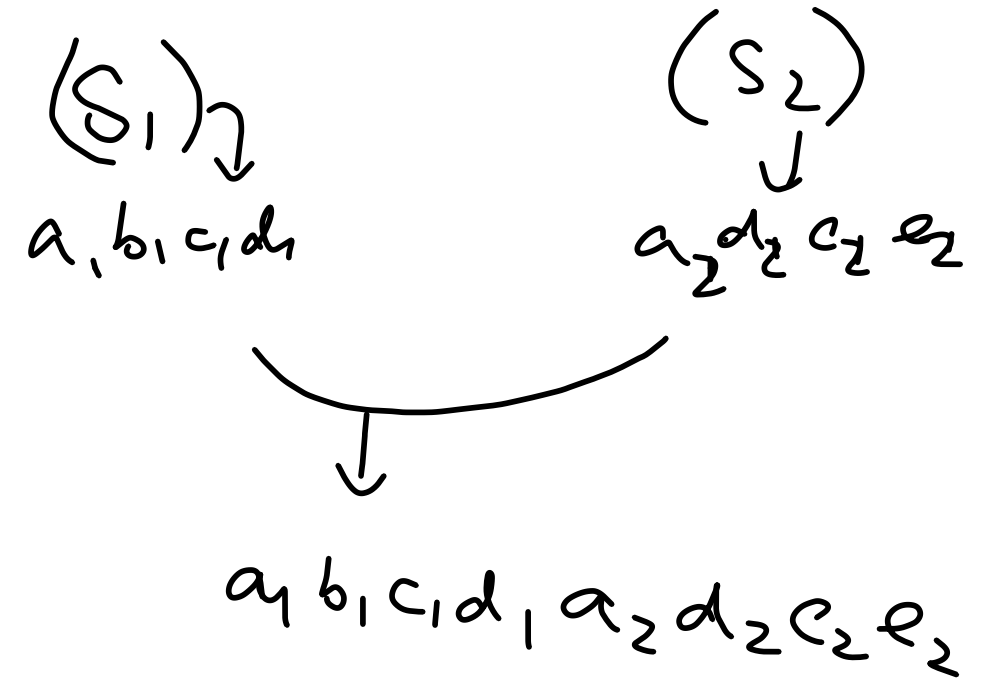
$$V = V_1 \cup V_2 \cup S$$

② Concat

$$P = P_1 \cup P_2 \cup \{ S \rightarrow S_1 S_2 \}$$

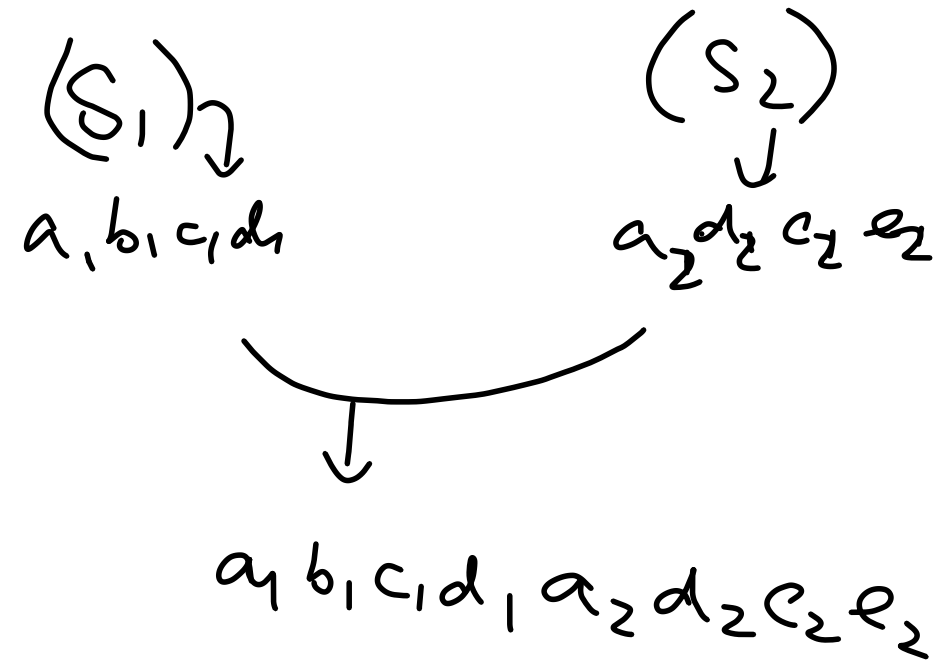
③ Star

(



② Concat

$$P = P_1 \cup P_2 \cup \{ S \rightarrow S_1 S_2 \}$$



③ Star

(V, Σ, P, S)

$$P = P_1 \cup \{ S \rightarrow S_1 S \mid \epsilon \}$$

④ Reversal

$$A \rightarrow A_1 A_2 \dots A_k$$

$$\Rightarrow A \rightarrow A_k \dots A_2 A_1$$

$$\left\{ \begin{array}{l} A \rightarrow \alpha B b \\ B \rightarrow 0 1 \mid 0 1 \end{array} \right.$$

$$\Downarrow$$

$$\left\{ \begin{array}{l} A \rightarrow b B a \\ B \rightarrow 0 1 \mid 1 0 \end{array} \right.$$

⑤ Intersection

$$L_1 = \{ a^m b^m c^n \mid m, n \geq 0 \} - \text{CFL}$$

$$L_2 = \{ a^m b^n c^n \mid m, n \geq 0 \} - \text{CFL}$$

$$L_1 \cap L_2 = \{ a^m b^m c^m \mid m, n \geq 0 \} - \underline{\text{(Not CFL)}}$$

$$\begin{aligned} S &\rightarrow S_1 S_2 \\ S_1 &\rightarrow a S_1 b \mid \epsilon \\ S_2 &\rightarrow c S_2 \mid \epsilon \end{aligned}$$

⑤ Intersection

$$L_1 = \{ a^m b^m c^n \mid m, n \geq 0 \} - \text{CFL}$$

$$L_2 = \{ a^m b^n c^n \mid m, n \geq 0 \} - \text{CFL}$$

$$L_1 \cap L_2 = \{ a^m b^m c^m \mid m, n \geq 0 \} - \text{(Not CFL)}$$

$$\begin{aligned} S &\rightarrow S_1 S_2 \\ S_1 &\rightarrow a S_1 b \mid \epsilon \\ S_2 &\rightarrow c S_2 \mid \epsilon \end{aligned}$$

⑤ Intersection

$$L_1 = \{ a^m b^m c^n \mid m, n \geq 0 \} - \text{CFL}$$

$$L_2 = \{ a^m b^n c^n \mid m, n \geq 0 \} - \text{CFL}$$

$$L_1 \cap L_2 = \{ a^m b^m c^m \mid m, n \geq 0 \} - (\text{Not CFL})$$

not closure

$$\begin{aligned} S &\rightarrow S_1 S_2 \\ S_1 &\rightarrow a S_1 b \mid \epsilon \\ S_2 &\rightarrow c S_2 \mid \epsilon \end{aligned}$$

⑥ Complement

$$L_1 \cap L_2 = \overline{(\overline{L_1} \cup \overline{L_2})}$$

Assume, $L \rightarrow \text{CFL} \Rightarrow \overline{L} \rightarrow \text{CFL}$

$$L_1, L_2 \rightarrow \text{CFL} \Rightarrow (L_1 \cap L_2 \rightarrow \text{CFL})$$

⑦ Set Difference

$$L_1 - L_2$$

$$= L_1 \cap \overline{L_2}$$

$$\overline{L} = \bigcup_{L \in \text{CFL}}^* \overline{L}$$

$$L = \{a^n b^n c^n \mid n \geq 0\}$$

$$\overline{L} = ? \downarrow$$

Thm

$L_1 \rightarrow \text{CFL}$

$L_2 \rightarrow \text{Regular}$

$L_1 \cap L_2 \rightarrow \text{CFL}$

$L = \{ \omega \mid \#a = \#b = \#c \}$

is not CFL.

$L_1 = a^* b^* c^* \Rightarrow \underline{\text{Regular}}$

$L \cap L_1 = \{ a^n b^n c^n \mid n \geq 0 \}$

$\Rightarrow L$ is not CFL

