

Runway Reservation Problem

- Single runway in an airport.
- Authority maintains landing times of each flight.
- When a new request comes with time t , checks if there is any landing reserved between $(t-3)$ to $(t+3)$. If not approve. O/w don't approve.

new_request(t)

if (- - - -)

- Condition Checking
- Insert

O/W

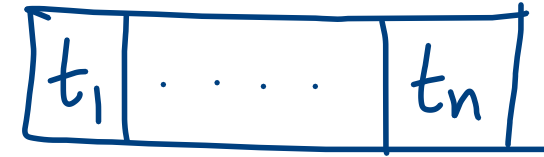
Approve

Don't Approve

$T = \{t_1, \dots, t_n\}$

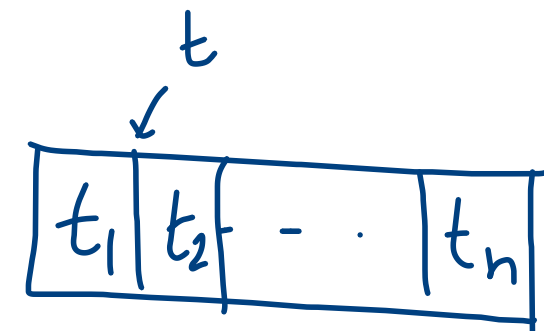
① Array :

(unsorted)



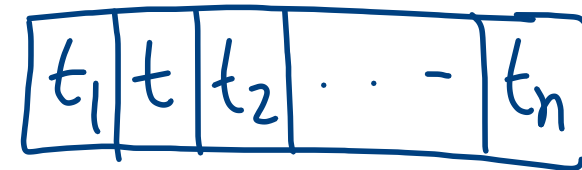
$O(n)$

(sorted)



Binary Search
 $\hookrightarrow O(\log n)$

$\Downarrow$



Insertion
 $\rightarrow O(n)$

② Linked List

$O(n)$

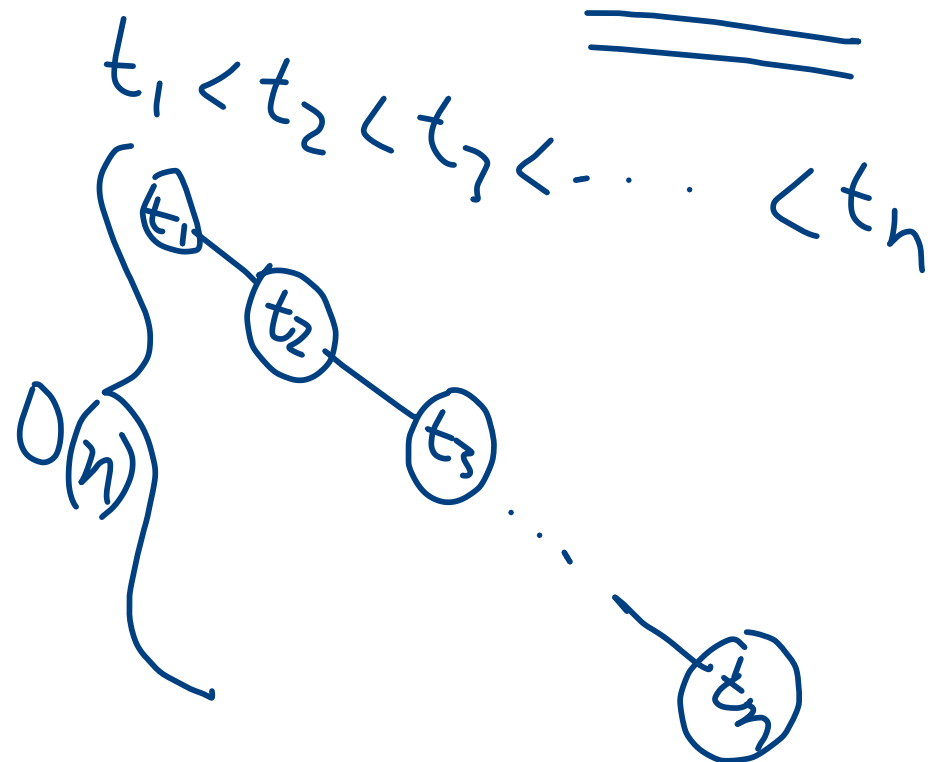
③

Search Binary Tree

(t_1)

Searching in BST

$\hookrightarrow O(h)$ or $O(\log n)$?



Condition $\rightarrow$ Searching in BST
 $\Rightarrow O(h)$

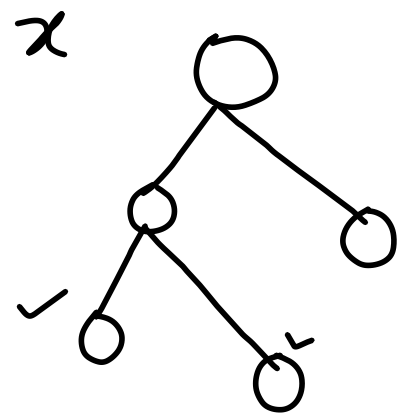
Insert $\rightarrow$ Insertion in BST
 $\Rightarrow O(h)$

Certain instances
 $h \approx O(n)$

Balanced Binary Search Tree

$$\underline{\underline{h = O(\log n)}}$$

AVL Tree (Adelson, Velsky, Landis)



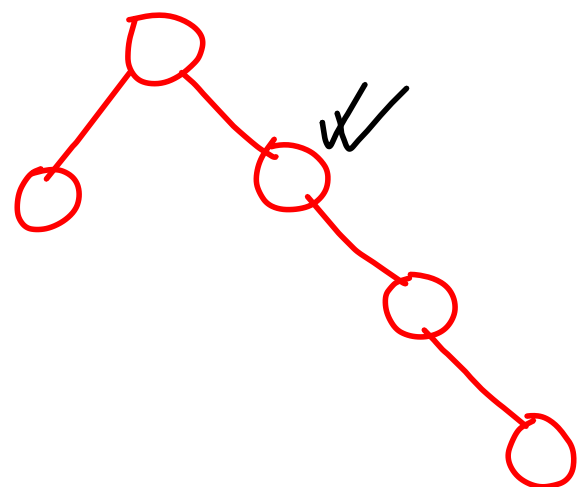
$$\text{height}(x) = 2$$

Height of node $\rightarrow$ max distance from that node to any leaf.

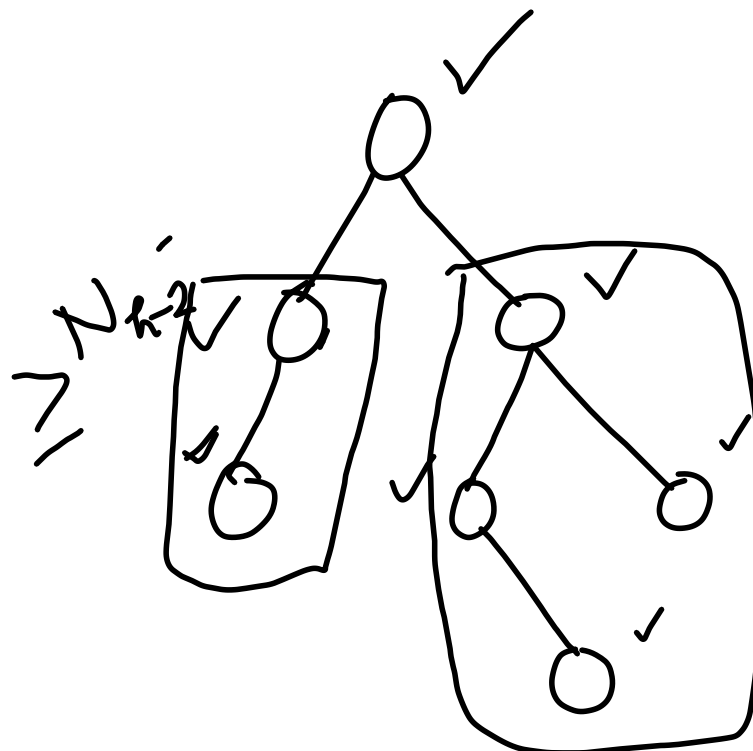
$\forall$ node x ,

AVL:

$$\forall x, \text{height}(x) = \max \{ \text{height}(lc(x)), \text{height}(rc(x)) \} + 1$$
$$| \text{height}(lc(x)) - \text{height}(rc(x)) | \leq 1$$



Not AVL



AVL $\geq N_{h-2}$

$$h = O(\log n)$$

Height of an AVL tree
is $O(\log n)$

$N_h =$ # nodes possible
with an AVL tree with h .

$$N_h \geq 2N_{h-2} + 1$$

$$\geq 2(2N_{h-4} + 1) + 1$$

$$= 2^2 N_{h-4} + (1+2)$$

...

$$= 2^i N_{h-2i} + (1+2+\dots+2^{i-1})$$

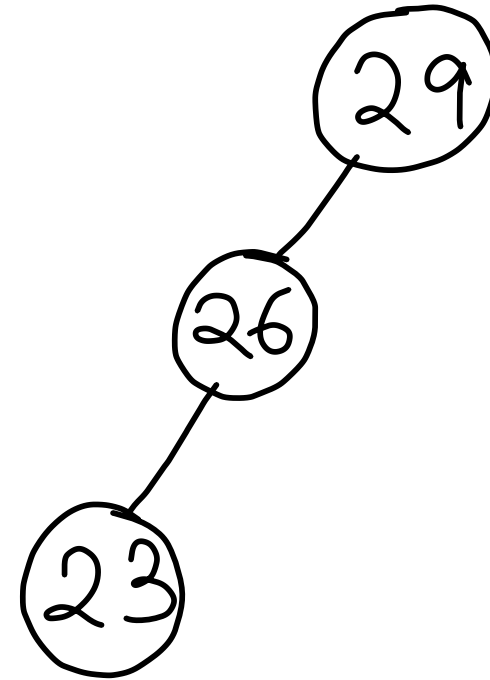
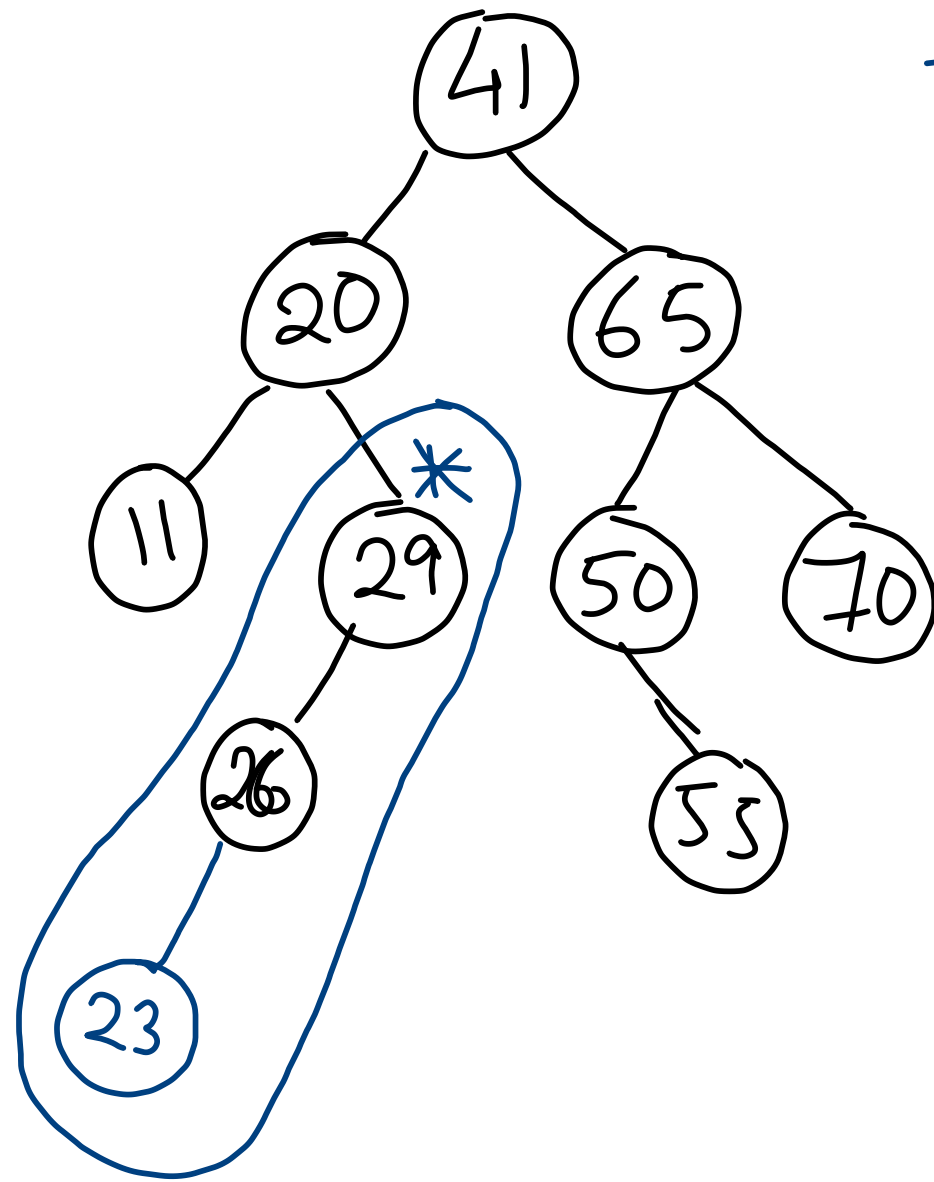
$(2^i - 1)$

$$N_h \geq 2^{\frac{h}{2}} N_0 + 2^{\frac{h}{2}}$$

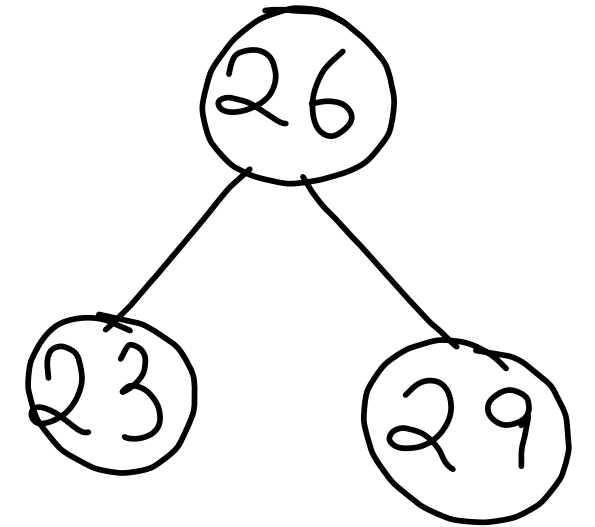
$$n \geq 2^{\frac{h}{2} + 1}$$

$$h = \log(n)$$

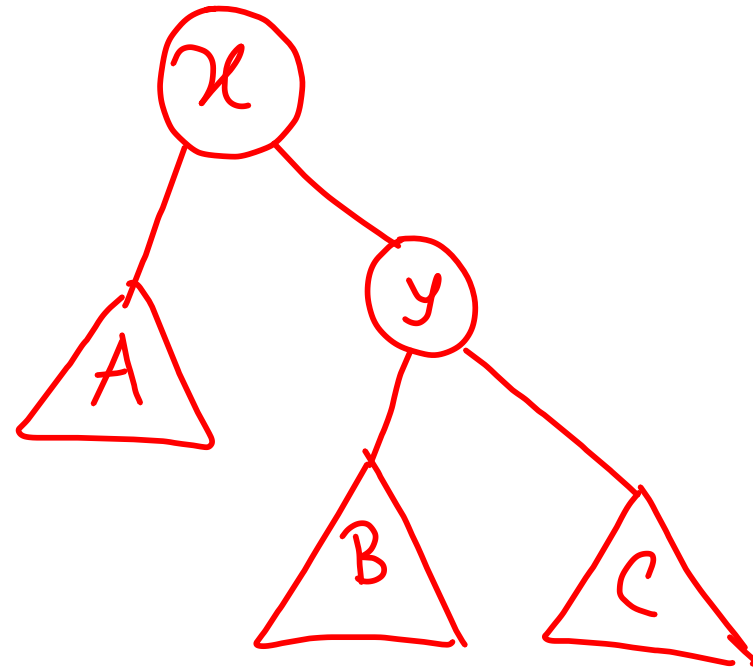
Insert(23)



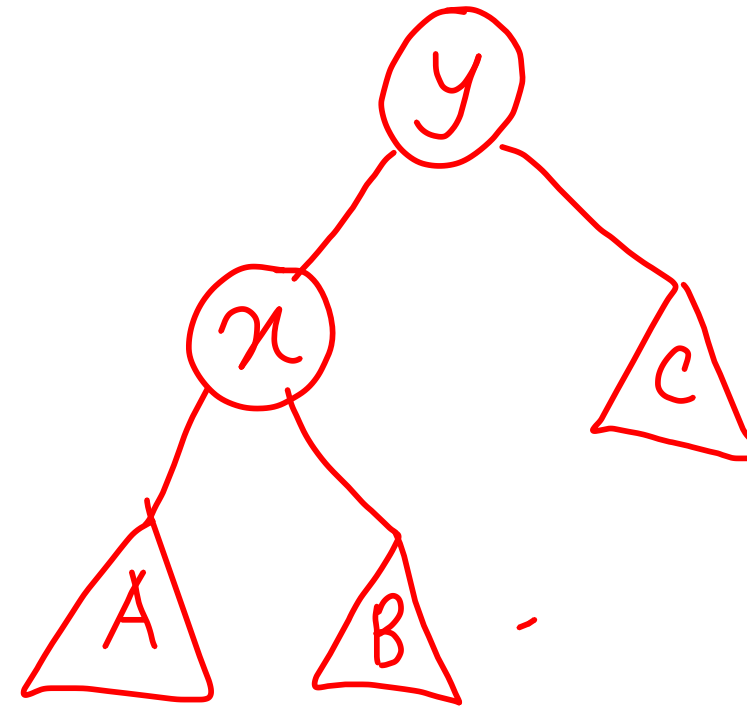
(Rotation)
⇒



Rotation



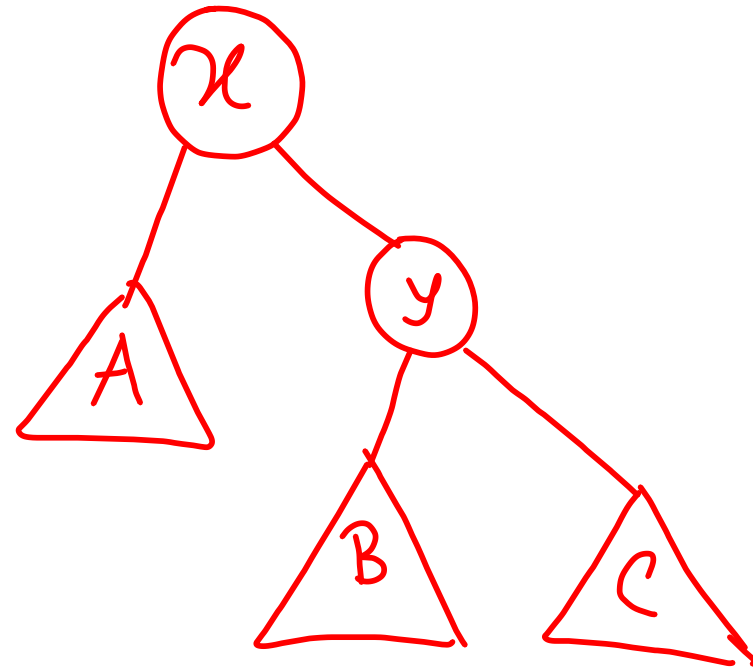
Left
Rotation
→
on
x



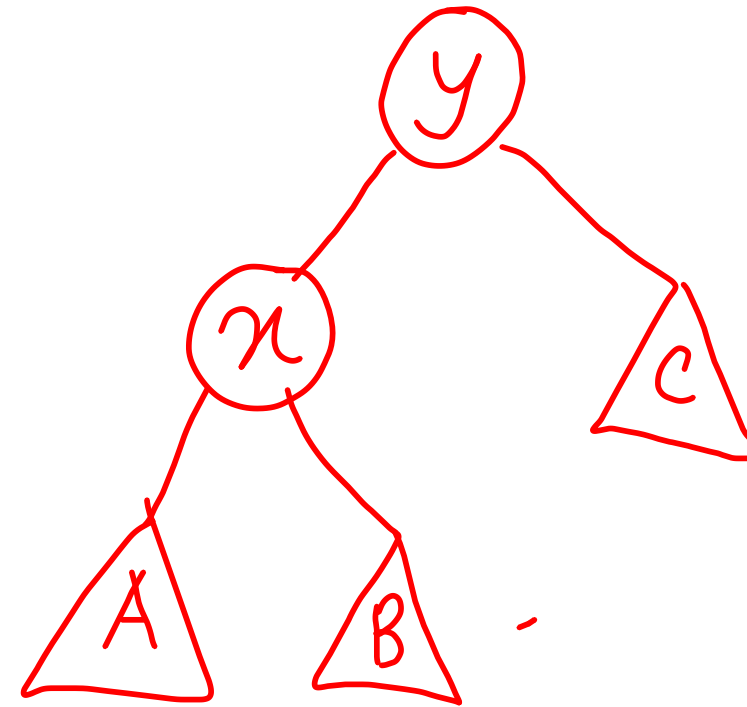
$A < x < B < y < C$

$A < x < B < y < C$

Rotation



Left
Rotation
→
on
x



$A < x < B < y < C$

$A < x < B < y < C$

AVL Insertion

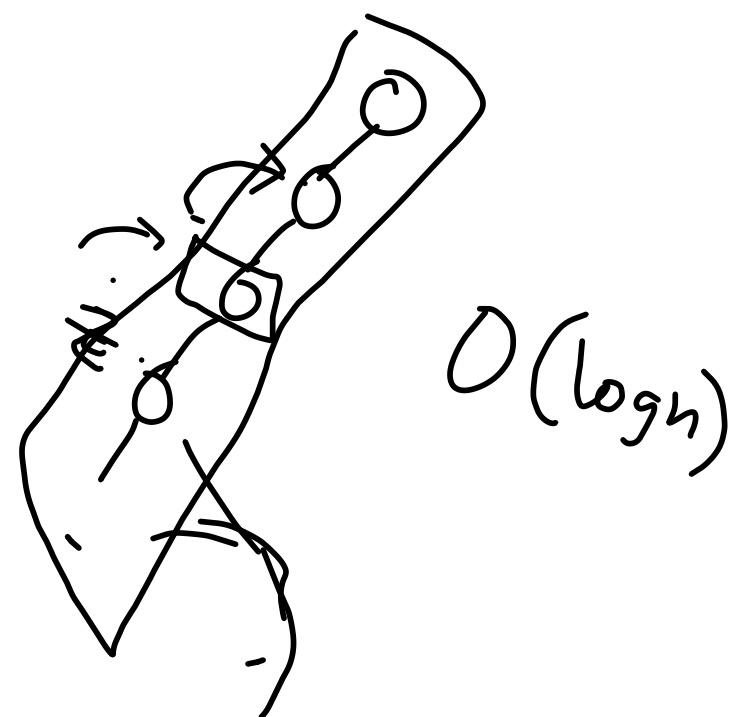
$$\underline{\underline{O(\log n)}}$$

- Insert w

- First element in the path from w to root which is unbalanced

Let's call that element $\textcircled{z}$

- Claim: z has a child & a grand-child in the path from z to w

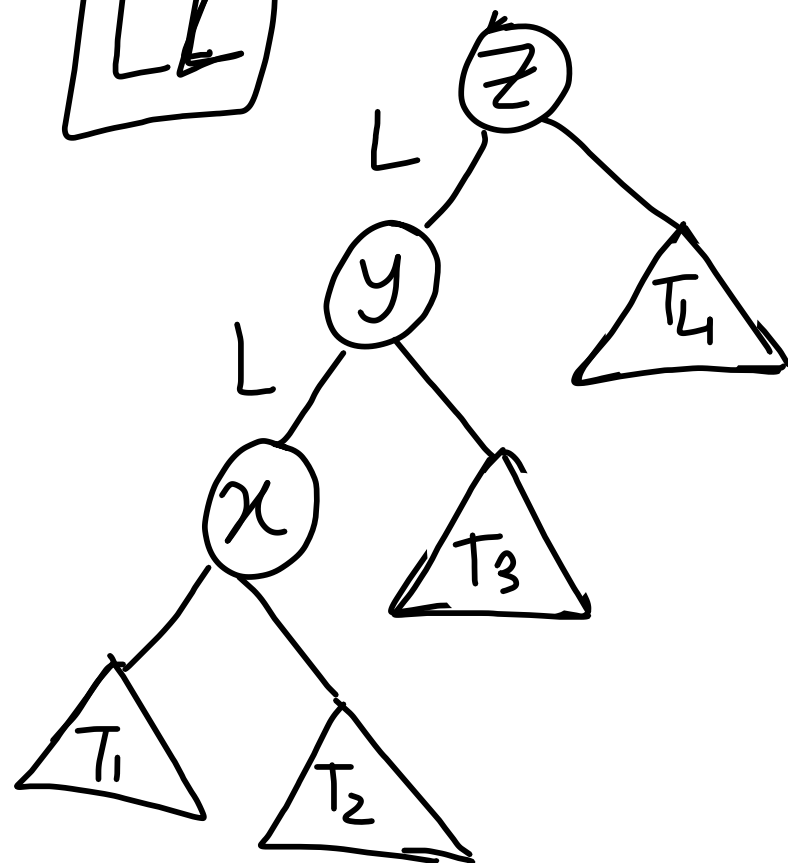


$\Downarrow$
y

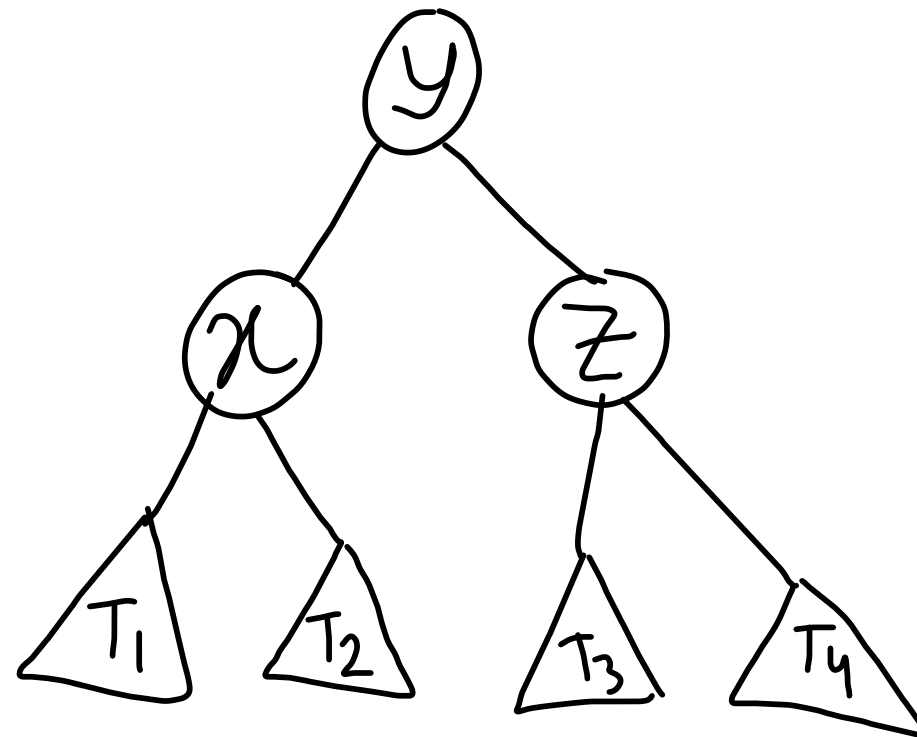
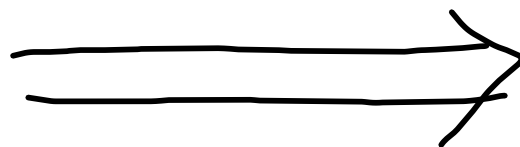
$\Downarrow$
x

(may include w itself)

LL

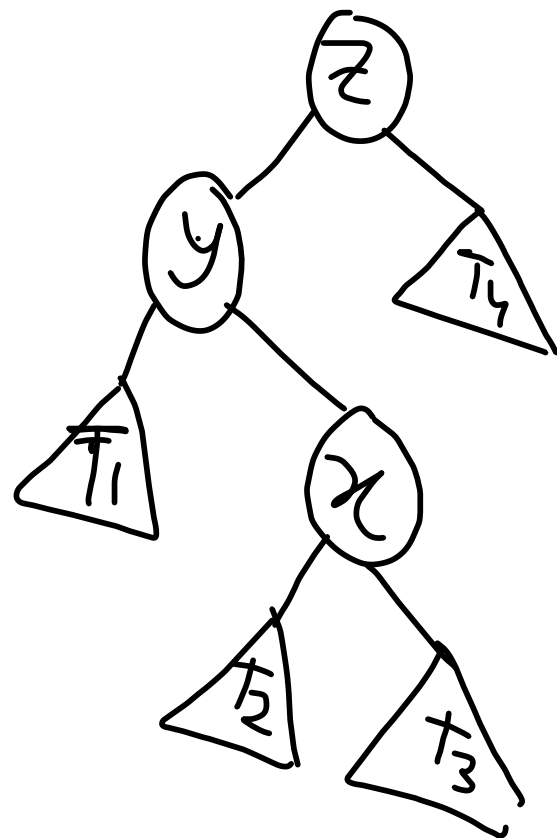


RR(z)

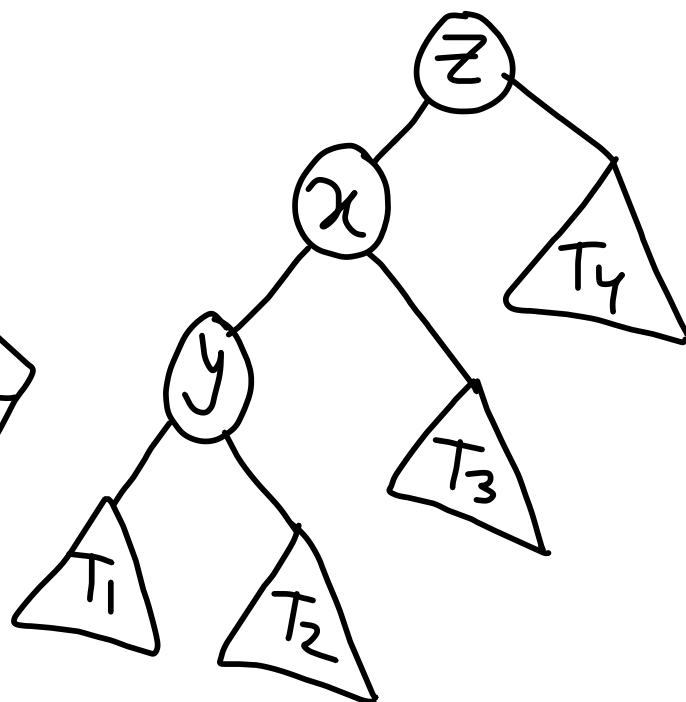


$O(\log n)$

LR



LR(y)



RR(z)

